

15. Recursion 2

Building a Calculator, Formal Grammars, Extended Backus Naur Form (EBNF), Parsing Expressions

Motivation: Calculator

Example

Input: $3 + 5$

Output: 8

- binary Operators $+$, $-$, $*$, $/$ and numbers

Motivation: Calculator

Example

Input: 3 / 5

Output: 0.6

- binary Operators +, -, *, / and numbers
- floating point arithmetic

Motivation: Calculator

Example

Input: $3 + 5 * 20$

Output: 103

- binary Operators $+$, $-$, $*$, $/$ and numbers
- floating point arithmetic
- precedences and associativities like in C++

Motivation: Calculator

Example

Input: $(3 + 5) * 20$

Output: 160

- binary Operators $+$, $-$, $*$, $/$ and numbers
- floating point arithmetic
- precedences and associativities like in C++
- parentheses

Motivation: Calculator

Example

Input: $-(3 + 5) + 20$

Output: 12

- binary Operators $+$, $-$, $*$, $/$ and numbers
- floating point arithmetic
- precedences and associativities like in C++
- parentheses
- unary operator $-$

Naive Attempt (without Parentheses)

```
double lval;
std::cin >> lval;

char op;
while (std::cin >> op && op != '=') {
    double rval;
    std::cin >> rval;

    if (op == '+')
        lval += rval;
    else if (op == '*')
        lval *= rval;
    else ...
}
std::cout << "Ergebnis " << lval << "\n";
```

Seems to work...

```
double lval;
std::cin >> lval;

char op;
while (std::cin >> op && op != '=') {
    double rval;
    std::cin >> rval;

    if (op == '+')
        lval += rval;
    else if (op == '*')
        lval *= rval;
    else ...
}

std::cout << "Ergebnis " << lval << "\n";
```

```
Input 1 * 2 * 3 * 4 =
Result 24
```


Oops, Multiplication first...

```
double lval;
std::cin >> lval;

char op;
while (std::cin >> op && op != '=') {
    double rval;
    std::cin >> rval;

    if (op == '+')
        lval += rval;
    else if (op == '*')
        lval *= rval;
    else ...
}
std::cout << "Ergebnis " << lval << "\n";
```

Input 2 + 3 * 3 =
Result 15

Analyzing the Problem

Example

Input:

$$13 + \dots$$

Analyzing the Problem

Example

Input:

$$13 + 4 * \dots$$

Analyzing the Problem

Example

Input:

$$13 + 4 * (15 - ...$$

Analyzing the Problem

Example

Input:

$$13 + 4 * (15 - 7 * ...$$

Analyzing the Problem

Example

Input:

$$13 + 4 * (15 - 7 * 3) =$$

Needs to be stored such that
evaluation can be performed

Analyzing the Problem

Example

Result:

$$13 + 4*(15 - 21)$$

Analyzing the Problem

Example

Result:

$$13 + 4 * (-6)$$

Analyzing the Problem

Example

Result:

$$13 + (-24)$$

Analyzing the Problem

Example

Result:

-11

Analyzing the Problem

Example

Expression:

$$13 + 4 * (15 - 7 * 3)$$

Example

This

Analyzing the Problem

Example

Expression:

$$13 + 4 * (15 - 7 * 3)$$

Example

This lecture

Analyzing the Problem

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Expression:

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This lecture is

Analyzing the Problem

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Expression:

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This lecture is pretty

Analyzing the Problem

Example

Expression:

$$13 + 4 * (15 - 7 * 3)$$

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This lecture is pretty much

Analyzing the Problem

Example

Expression:

$$13 + 4 * (15 - 7 * 3)$$

Example

This lecture is pretty much recursive.

Analyzing the Problem

$$13 + 4 * (15 - 7 * 3)$$

“Understanding an expression requires lookahead to upcoming symbols!

We will store symbols elegantly using recursion.

We need a new formal tool (that is independent of C++).

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Formal Grammars

- Alphabet: finite set of symbols
- Strings: finite sequences of symbols

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A formal grammar defines which strings are valid.

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- Alphabet: finite set of symbols
- Strings: finite sequences of symbols

A formal grammar defines which strings are valid.

To describe the formal grammar, we use:

Extended Backus Naur Form (EBNF)

What Can We Do about the Unnecessary Diversity of Notation for Syntactic Definitions?

Niklaus Wirth
Federal Institute of Technology (ETH), Zürich, and
Xerox Palo Alto Research Center

Key Words and Phrases: syntactic description
language, extended BNF
CR Categories: 4.20

The population of programming languages is steadily growing, and there is no end of this growth in sight. Many language definitions appear in journals, many are found in technical reports, and perhaps an even greater number remains confined to proprietary circles. After frequent exposure to these definitions, one cannot fail to notice the lack of "common denominators." The only widely accepted fact is that the language structure is defined by a syntax. But even notation for syntactic description eludes any commonly agreed standard form, although the underlying ancestor is invariably the Backus-Naur Form of the Algol 60 report. As variations are often only slight, they become annoying for their very lack of an apparent motivation.

Out of sympathy with the troubled reader who is weary of adapting to a new variant of BNF each time another language definition appears, and without any claim for originality, I venture to submit a simple notation that has proven valuable and satisfactory in use. It has the following properties to recommend it:

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Author's present address: Xerox Corporation, Palo Alto Research Center, 3333 Coyote Hill Road, Palo Alto, CA 94304.

1. The notation distinguishes clearly between meta-, terminal, and nonterminal symbols.
2. It does not exclude characters used as metasympols from use as symbols of the language (as e.g. "|" in BNF).
3. It contains an explicit iteration construct, and thereby avoids the heavy use of recursion for expressing simple repetition.
4. It avoids the use of an explicit symbol for the empty string (such as (empty) or ϵ).
5. It is based on the ASCII character set.

This meta language can therefore conveniently be used to define its own syntax, which may serve here as an example of its use. The word *identifier* is used to denote *nonterminal symbol*, and *literal* stands for *terminal symbol*. For brevity, *identifier* and *character* are not defined in further detail.

```
syntax      = {production}.
production = identifier "=" expression " ".
expression = term {"|" term}.
term       = factor {factor}.
factor     = identifier | literal | "(" expression ")" |
            "[" expression "]" | "{" expression "}".
literal    = " " " " character {character} " " " " .
```

Repetition is denoted by curly brackets, i.e. {a} stands for ϵ | a | aa | aaa | Optionality is expressed by square brackets, i.e. [a] stands for ϵ | a. Parentheses merely serve for grouping, e.g. (a|b|c stands for ac | bc. Terminal symbols, i.e. literals, are enclosed in quote marks (and, if a quote mark appears as a literal itself, it is written twice), which is consistent with common practice in programming languages.

Received January 1977; revised February 1977

Expressions

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

Expressions

$$- (\underline{3} - (\underline{4} - \underline{5})) * (\underline{3} + \underline{4} * \underline{5}) / \underline{6}$$

What do we need in a grammar?

- Number

Expressions

$$\underline{-} \left(\underline{3} - \underline{(4 - 5)} \right) * \left(\underline{3} + \underline{4 * 5} \right) / \underline{6}$$

What do we need in a grammar?

- Number , (?)

Expressions

$$\underline{-} (3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)

Expressions

$$-(3 - (4 - 5)) \underline{*} (3 + \underline{4} \underline{*} 5) \underline{/} 6$$

What do we need in a grammar?

- Number, (?)
-Number, -(?)
- ? * ?, ? / ?, ...

Expressions

$$-(3_{\underline{\quad}} - (4_{\underline{\quad}} - 5)) * (3_{\underline{\quad}} + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- ? * ?, ? / ?, ...
- ? - ?, ? + ?, ...

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- ? * ?, ? / ?, ...
- ? - ?, ? + ?, ...

Factor

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- Factor * Factor,
Factor / Factor , ...
- ? - ?, ? + ?, ...

Factor

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- Factor * Factor,
Factor / Factor , ...
- ? - ?, ? + ?, ...

Factor

Term

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- Factor * Factor, Factor
Factor / Factor , ...
- ? - ?, ? + ?, ...

Factor

Term

Expressions

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- Factor * Factor, Factor
Factor / Factor , ...
- Term + Term,
Term - Term, ...

Factor

Term

Expressions

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- Factor * Factor, Factor
Factor / Factor , ...
- Term + Term,
Term - Term, ...

Factor

Term

Expression

Expressions

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (?)
-Number, -(?)
- Factor * Factor, Factor
Factor / Factor , ...
- Term + Term, **Term**
Term - Term, ...

Factor

Term

Expression

Expressions

$$-(3 - (4 - 5)) * (3 + 4 * 5) / 6$$

What do we need in a grammar?

- Number , (Expression)
-Number, -(Expression)
- Factor * Factor, Factor
Factor / Factor , ...
- Term + Term, Term
Term - Term, ...

Factor

Term

Expression

The EBNF for Expressions

A factor is

- a number,
- an expression in parentheses or
- a negated factor.

```
factor      = number  
            | "(" expression ")"  
            | "-" factor.
```

The EBNF for Expressions

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The EBNF for Expressions

A factor is

- a number,
- an expression in parentheses or
- a negated factor.

non-terminal symbol

factor = number
| "(" expression ")"
| "-" factor.

terminal symbol

alternative

The EBNF for Expressions

A term is

- factor,
- factor * factor, factor / factor,
- factor * factor * factor, factor / factor * factor, ...
- ...

term = factor { "*" factor | "/" factor }.

The EBNF for Expressions

A term is

- factor,
- factor * factor, factor / factor,
- factor * factor * factor, factor / factor * factor, ...
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term = factor { "*" factor | "/" factor }.

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The EBNF for Expressions

A term is

- factor,
- factor * factor, factor / factor,
- factor * factor * factor, factor / factor * factor, ...
- ...

term = factor { "*" factor | "/" factor }.

The EBNF for Expressions

A term is

- factor,
- factor * factor, factor / factor,
- factor * factor * factor, factor / factor * factor, ...
- ...

term = factor { "*" factor | "/" factor } .

optional repetition

The EBNF for Expressions

factor = number
| "(" expression ")"
| "-" factor.

term = factor { "*" factor | "/" factor }.

expression = term { "+" term | "-" term }.

Numbers

An *integer* comprises at least one digit, followed by an arbitrary number of digits.

```
number = digit { digit }.  
digit  = '0' | '1' | '2' | ... | '9'.
```

Numbers

An integer *comprises at least one digit*, followed by an arbitrary number of digits.

```
number = digit { digit }.  
digit  = '0' | '1' | '2' | ... | '9'.
```

Numbers

An integer comprises at least one digit, *followed by an arbitrary number of digits.*

number = digit { digit }.

digit = '0' | '1' | '2' | ... | '9'.

Parsing

- **Parsing:** Check if a string is valid according to the EBNF.

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- **Parser:** A program for parsing.

Parsing

- **Parsing:** Check if a string is valid according to the EBNF.
- **Parser:** A program for parsing.
- **Useful:** From the EBNF we can (nearly) automatically generate a parser

Construct a Parser

- Rules become functions
- Alternatives and options become `if`-statements.
- Nonterminal symbols on the right hand side become function calls
- Optional repetitions become `while`-statements

Rules (except number)

factor = number
| "(" expression ")"
| "-" factor.

term = factor { "*" factor | "/" factor }.

expression = term { "+" term | "-" term }.

Expression is read from an input stream.

```
// POST: returns true if and only if is = factor ...  
//       and in this case extracts factor from is  
bool factor (std::istream& is);
```

```
// POST: returns true if and only if is = term ...,  
//       and in this case extracts all factors from is  
bool term (std::istream& is);
```

```
// POST: returns true if and only if is = expression ...,  
//       and in this case extracts all terms from is  
bool expression (std::istream& is);
```

Expression is read from an input stream.

```
// POST: extracts a factor from is  
//       and returns its value  
double factor (std::istream& is);
```

```
// POST: extracts a term from is  
//       and returns its value  
double term (std::istream& is);
```

```
// POST: extracts an expression from is  
//       and returns its value  
double expression (std::istream& is);
```

One Character Lookahead...

... to find the right alternative.

```
// POST: leading whitespace characters are extracted
//       from input, and the first non-whitespace character
//       input returned (0 if there input no such character)
char lookahead (std::istream& input)
{
    input >> std::ws;           // skip whitespaces
    if (input.eof())
        return 0;             // end of stream
    else
        return input.peek();   // next character in input
}
```

Cherry-Picking

... to extract the desired character.

```
// POST: if ch matches the next lookahead then consume it and return true
//      otherwise return false
bool consume (std::istream& input, char c)
{
    if (lookahead (input) == c) {
        input >> c;
        return true;
    } else
        return false ;
}
```

Evaluating Factors

```
double factor (std::istream& input)
{
    double value;
    if (consume (input, '(')) {
        value = expression (input);           // "(" expression
        consume (input, ')');                // ")"
    } else if (consume (input, '-'))
        value = -factor (input);           // - factor
    else
        value = number(input);
    return value;
}
```

```
factor = "(" expression ")"
        | "-" factor
        | number.
```

Evaluating Terms

```
double term (std::istream& input)
{
    double value = factor (input);           // factor
    while (true) {
        if (consume (input, '*'))           // "*" factor
            value *= factor (input);
        else if (consume (input, '/'))      // "/" factor
            value /= factor (input);
        else
            return value;
    }
}
```

term = factor { "*" factor | "/" factor }.

Evaluating Expressions

```
double expression (std::istream& input)
{
    double value = term (input);           // term
    while (true) {
        if (consume (input, '+'))
            value += term (input);         // "+" term
        else if (consume (input, '-'))
            value -= term (input);         // "-" term
        else
            return value;
    }
}
```

expression = term { "+" term | "-" term }.

Digits ...

```
// POST: returns the digit that could be consumed from a stream
//      (0 if no digit available)
// digit = '0' | '1' | ... | '9'.
char digit(std::istream& input){
    char ch = input.peek(); // one symbol lookahead
    if (input.eof()) return 0; // nothing available on the stream
    if (ch >= '0' && ch <= '9'){
        input >> ch; // consume
        return ch;
    }
    return 0;
}
```

... and Numbers

```
// POST: returns an unsigned integer consumed from the stream
// number = digit {digit}.
unsigned int number (std::istream& input){
    input >> std::skipws; // skip whitespaces before the first digit
    char ch = digit(input);
    input >> std::noskipws; // no whitespaces allowed within a number
    unsigned int num = 0;
    while(ch > 0){ // skip remaining digits
        num = num * 10 + ch - '0';
        ch = digit(input);
    }
    return num;
}
```

Recursion!

number

factor

term

expression

Recursion!

number

factor

term

expression

```
graph BT; expression --> term; factor; number;
```

Recursion!

number

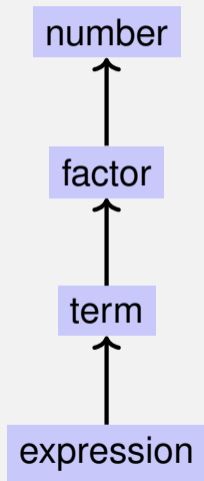
factor

term

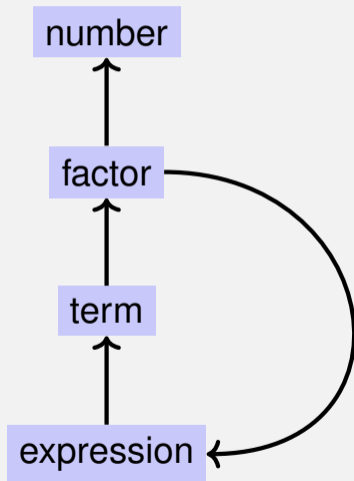
expression



Recursion!



Recursion!



EBNF — and it works!

EBNF (calculator.cpp, Evaluation from left to right):

```
factor    = number  
          | "(" expression ")"  
          | "-" factor.
```

```
term      = factor { "*" factor | "/" factor }.
```

```
expression = term { "+" term | "-" term }.
```

```
std::stringstream input ("1-2-3");  
std::cout << expression (input) << "\n"; // -4
```


16. Structs

Rational Numbers, Struct Definition, Function- and Operator Overloading

Calculating with Rational Numbers

- Rational numbers (\mathbb{Q}) are of the form $\frac{n}{d}$ with n and d in \mathbb{Z}
- C++ does not provide a built-in type for rational numbers

Calculating with Rational Numbers

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- C++ does not provide a built-in type for rational numbers

Goal

We build a C++-type for rational numbers ourselves! 😊

Vision

```
// input
std::cout << "Rational number r =? ";
rational r;
std::cin >> r;
std::cout << "Rational number s =? ";
rational s;
std::cin >> s;

// computation and output
std::cout << "Sum is " << r + s << ".\n";
```

A First Struct

```
struct rational {  
    int n;  
    int d; // INV: d != 0  
};
```

A First Struct

```
struct rational {  
    int n; ← member variable (numerator)  
    int d; ← // INV: d != 0  
};  
           ← member variable (denominator)
```

A First Struct

```
struct rational {  
    int n; ← member variable  
    int d; // INV: d != 0  
};  
           ← member variable
```

- struct defines a new *type*

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```

- struct defines a new *type*
- formal range of values: *cartesian product* of the value ranges of existing types

A First Struct

```
struct rational {  
    int n; ← member variable  
    int d; // INV: d != 0  
};  
           ← member variable
```

- struct defines a new *type*
- formal range of values: *cartesian product* of the value ranges of existing types
- real range of values: `rational` \subsetneq `int` \times `int`.

Accessing Member Variables

```
struct rational {
    int n;
    int d; // INV: d != 0
};

rational add (rational a, rational b){
    rational result;
    result.n = a.n * b.d + a.d * b.n;
    result.d = a.d * b.d;
    return result;
}
```

$$\frac{r_n}{r_d} := \frac{a_n}{a_d} + \frac{b_n}{b_d} = \frac{a_n \cdot b_d + a_d \cdot b_n}{a_d \cdot b_d}$$

Input

```
// Input r
rational r;
std::cout << "Rational number r:\n";
std::cout << " numerator =? ";
std::cin >> r.n;
std::cout << " denominator =? ";
std::cin >> r.d;

// Input s the same way
rational s;
...
```

Vision comes within Reach ...

```
// computation
const rational t = add (r, s);

// output
std::cout << "Sum is " << t.n << "/" << t.d << ".\n";
```

Struct Definitions: Examples

```
struct rational_vector_3 {  
    rational x;  
    rational y;  
    rational z;  
};
```

underlying types can be fundamental or **user defined**

Struct Definitions: Examples

```
struct extended_int {  
    // represents value if is_positive==true  
    // and -value otherwise  
    unsigned int value;  
    bool is_positive;  
};
```

the underlying types can be **different**

Structs: Initialization and Assignment

```
rational s; ← member variables are uninitialized
```

```
rational t = {1,5};
```

```
rational u = t;
```

```
t = u;
```

```
rational v = add (u,t);
```

Structs: Initialization and Assignment

```
rational s;
```

```
rational t = {1,5}; ← member-wise initialization:  
t.n = 1, t.d = 5
```

```
rational u = t;
```

```
t = u;
```

```
rational v = add (u,t);
```


Structs: Initialization and Assignment

```
rational s;
```

```
rational t = {1,5};
```

```
rational u = t; ← member-wise copy
```

```
t = u;
```

```
rational v = add (u,t);
```

Structs: Initialization and Assignment

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rational s;
```

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rational t = {1,5};
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rational u = t;
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```
t = u; ← member-wise copy
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```
rational v = add (u,t);
```

Structs: Initialization and Assignment

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rational s;
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rational t = {1,5};
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t = u;
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rational v = add (u,t); ← member-wise copy
```

Comparing Structs?

For each fundamental type (`int`, `double`, ...) there are comparison operators `==` and `!=`, not so for structs! Why?

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- member-wise comparison does not make sense in general...

Comparing Structs?

For each fundamental type (`int`, `double`, ...) there are comparison operators `==` and `!=`, not so for structs! Why?

- member-wise comparison does not make sense in general...
- ...otherwise we had, for example, $\frac{2}{3} \neq \frac{4}{6}$

User Defined Operators

Instead of

```
rational t = add(r, s);
```

we would rather like to write

```
rational t = r + s;
```

User Defined Operators

Instead of

```
rational t = add(r, s);
```

we would rather like to write

```
rational t = r + s;
```

This can be done with *Operator Overloading*.

Function Overloading

- A function is defined by name, types, number and order of arguments

```
double sq (double x) { ... }           // f1
int sq (int x) { ... }                 // f2
int pow (int b, int e) { ... }        // f3
int pow (int e) { return pow (2,e); } // f4
```

- the compiler automatically chooses the function that fits “best” for a function call

```
std::cout << sq (3);
std::cout << sq (1.414);
std::cout << pow (2);
std::cout << pow (3,3);
```

Function Overloading

- A function is defined by name, types, number and order of arguments

```
double sq (double x) { ... }           // f1
int sq (int x) { ... }                 // f2
int pow (int b, int e) { ... }        // f3
int pow (int e) { return pow (2,e); } // f4
```

- the compiler automatically chooses the function that fits “best” for a function call

```
std::cout << sq (3);
std::cout << sq (1.414);
std::cout << pow (2);
std::cout << pow (3,3);
```

Function Overloading

- A function is defined by name, types, number and order of arguments

```
double sq (double x) { ... }           // f1
int sq (int x) { ... }                 // f2
int pow (int b, int e) { ... }        // f3
int pow (int e) { return pow (2,e); } // f4
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```

Operator Overloading

- Operators are special functions and can be overloaded
- Name of the operator *op*:

```
operatorop
```


Adding rational Numbers – Before

```
// POST: return value is the sum of a and b
rational add (rational a, rational b)
{
    rational result;
    result.n = a.n * b.d + a.d * b.n;
    result.d = a.d * b.d;
    return result;
}

...
const rational t = add (r, s);
```

Adding rational Numbers – After

```
// POST: return value is the sum of a and b
rational operator+ (rational a, rational b)
{
    rational result;
    result.n = a.n * b.d + a.d * b.n;
    result.d = a.d * b.d;
    return result;
}
...
const rational t = r + s;
```

Adding rational Numbers – After

```
// POST: return value is the sum of a and b
rational operator+ (rational a, rational b)
{
    rational result;
    result.n = a.n * b.d + a.d * b.n;
    result.d = a.d * b.d;
    return result;
}
...
const rational t = r + s;
```

↑
infix notation

Adding rational Numbers – After

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rational operator+ (rational a, rational b)
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    rational result;
    result.n = a.n * b.d + a.d * b.n;
    result.d = a.d * b.d;
    return result;
}
...
const rational t = operator+ (r, s);
```

↑
equivalent but less handy: functional notation

Unary Minus

Only one argument:

```
// POST: return value is  $-a$   
rational operator- (rational a)  
{  
    a.n =  $-a.n$ ;  
    return a;  
}
```

Comparison Operators

can be defined such that they do the right thing:

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```
// POST: returns true iff a == b
bool operator==(rational a, rational b)
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    return a.n * b.d == a.d * b.n;
}
```

Comparison Operators

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}
```

$$\frac{2}{3} = \frac{4}{6} \quad \checkmark$$

Arithmetic Assignment

We want to write

```
rational r;  
r.n = 1; r.d = 2;           // 1/2
```

```
rational s;  
s.n = 1; s.d = 3;         // 1/3
```

```
r += s;  
std::cout << r.n << "/" << r.d;   // 5/6
```

Operator +=

```
rational& operator+= (rational& a, rational b)
{
    a.n = a.n * b.d + a.d * b.n;
    a.d *= b.d;
    return a;
}
```

Operator +=

```
rational& operator+= (rational& a, rational b)
{
    a.n = a.n * b.d + a.d * b.n;
    a.d *= b.d;
    return a;
}
```

- The L-value a is increased by the value of b and returned as L-value

In/Output Operators

can also be overloaded.

■ Before:

```
std::cout << "Sum is "  
          << t.n << "/" << t.d << "\n";
```

■ After (desired):

```
std::cout << "Sum is "  
          << t << "\n";
```

In/Output Operators

can be overloaded as well:

```
// POST: r has been written to out
std::ostream& operator<< (std::ostream& out,
                          rational r)
{
    return out << r.n << "/" << r.d;
}
```

In/Output Operators

can be overloaded as well:

```
// POST: r has been written to out
std::ostream& operator<< (std::ostream& out,
                          rational r)
{
    return out << r.n << "/" << r.d;
}
```

writes `r` to the output stream
and returns the stream as L-value.

Input

```
// PRE: in starts with a rational number
// of the form "n/d"
// POST: r has been read from in
std::istream& operator>> (std::istream& in,
                          rational& r){
    char c; // separating character '/'
    return in >> r.n >> c >> r.d;
}
```

reads `r` from the input stream
and returns the stream as L-value.

Goal Attained!

```
// input
std::cout << "Rational number r =? ";
rational r;
std::cin >> r;

std::cout << "Rational number s =? ";
rational s;
std::cin >> s;

// computation and output
std::cout << "Sum is " << r + s << ".\n";
```


Goal Attained!

```
// input
std::cout << "Rational number r =? ";
rational r;
std::cin >> r;

std::cout << "Rational number s =? ";
rational s;
std::cin >> s;

// computation and output
std::cout << "Sum is " << r + s << ".\n";
```

operator >>

operator +

operator <<