

# Datenstrukturen und Algorithmen

## Exercise 6

FS 2021

# Program of today

- 1 Feedback of last exercise
- 2 Repetition theory
  - Binary Trees
- 3 Repetition Theory
  - AVL Condition
  - AVL Insert

# Feedback

Open hashing:

- $h'(k) = \lceil \ln(k + 1) \rceil \bmod q$

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- $s(j, k) = ((k \cdot j) \bmod q) + 1$

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Open hashing:

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- $s(j, k) = k^j \bmod p \rightarrow$  not suitable:  $(k = 0) \mapsto 0, (k = 1) \mapsto 1$
- $s(j, k) = ((k \cdot j) \bmod q) + 1 \rightarrow$  not suitable: 1 if  $k$  is multiple of  $q$ , and range  $p - q$  is not covered



# Feedback

## Coocoo hashing

- $h_1(k) = k \bmod 5$ ,  $h_2(k) = \lfloor k/5 \rfloor \bmod 5$
- add 27, 2, 32

T\_1: \_\_, \_\_, 27, \_\_, \_\_

T\_2: \_\_, \_\_, \_\_, \_\_, \_\_

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T\_1: \_\_, \_\_, 27, \_\_, \_\_

T\_2: 2, 32, \_\_, \_\_, \_\_

# Feedback

## Coocoo hashing

- $h_1(k) = k \bmod 5$ ,  $h_2(k) = \lfloor k/5 \rfloor \bmod 5$
- add 7: infinite loop

|     |      |    |   |    |   |    |   |    |   |    |  |      |    |   |    |   |    |   |    |   |    |  |
|-----|------|----|---|----|---|----|---|----|---|----|--|------|----|---|----|---|----|---|----|---|----|--|
|     | T_1: | __ | , | __ | , | 27 | , | __ | , | __ |  | T_2: | 2  | , | 32 | , | __ | , | __ | , | __ |  |
| 7:  | T_1: | __ | , | __ | , | 7  | , | __ | , | __ |  | T_2: | 27 | , | 32 | , | __ | , | __ | , | __ |  |
| 2:  | T_1: | __ | , | __ | , | 2  | , | __ | , | __ |  | T_2: | 27 | , | 7  | , | __ | , | __ | , | __ |  |
| 32: | T_1: | __ | , | __ | , | 32 | , | __ | , | __ |  | T_2: | 2  | , | 7  | , | __ | , | __ | , | __ |  |
| 27: | T_1: | __ | , | __ | , | 27 | , | __ | , | __ |  | T_2: | 2  | , | 32 | , | __ | , | __ | , | __ |  |
| 7:  | ...  |    |   |    |   |    |   |    |   |    |  |      |    |   |    |   |    |   |    |   |    |  |

# Feedback

## Finding a Sub-Array

```
// calculating hash_a, hash_b, c_to_k
It1 window_end = from;
for(It2 current = begin; current != end;
    ++current, ++window_end) {
    if(window_end == to) return to;
    hash_b = (C * hash_b % M + *current) % M;
    hash_a = (C * hash_a % M + *window_end) % M;
    c_to_k = c_to_k * C % M;
}
```

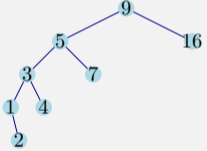
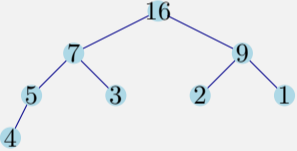
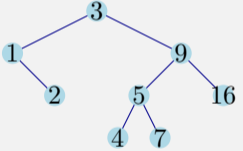
# Feedback

## Finding a Sub-Array

```
// looking for b and updating hash_a
for(It1 window_begin = from; ;
    ++window_begin, ++window_end) {
    if(hash_a == hash_b)
        if(std::equal(window_begin, window_end, begin, end))
            return window_begin;
    if(window_end == to) return to;
    hash_a = (C * hash_a % M + *window_end
              + (M - c_to_k) * *window_begin % M) % M;
}
```

## **2. Repetition theory**

# Comparison of binary Trees

|           | Search trees  | Heaps<br>Min- / Max-<br>Heap   | Balanced trees<br>AVL, red-black tree   |
|-----------|---|--|---|
| in C++:   |   | <code>std::make_heap</code>  | <code>std::map</code>   |
|           |  |  |  |
| Insertion | $\Theta(h(T))$  | $\Theta(\log n)$   | $\Theta(\log n)$  |
| Search    | $\Theta(h(T))$  | $\Theta(n)$ (!!)   | $\Theta(\log n)$  |
| Deletion  | $\Theta(h(T))$  | Search + $\Theta(\log n)$  | $\Theta(\log n)$  |

# Comparison of binary Trees

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|           |                |                              |                                       |
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| Search    | $\Theta(h(T))$ | $\Theta(n)$ (!!)             | $\Theta(\log n)$                      |
| Deletion  | $\Theta(h(T))$ | Search + $\Theta(\log n)$    | $\Theta(\log n)$                      |

**Recall:**  $\Theta(\log n) \leq \Theta(h(T)) \leq \Theta(n)$

# Repetition: Binary Trees, Inserting a Key

## Binary Search Trees

- Search for Key.
- Insert at the reached empty leaf (`null`).

## MinHeap

- Insert at the very back of the Array.
- Restore Heap-Condition: `siftUp` (climb successively).



# Repetition: Binary Trees, Inserting a Key

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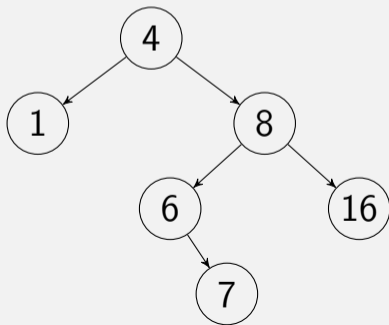
- Insert at the very back of the Array.
- Restore Heap-Condition: `siftUp` (climb successively).

**Exercise:** Insert 4, 8, 16, 1, 6, 7 into empty Tree/Heap.

# Repetition: Binary Trees, Inserting a Key

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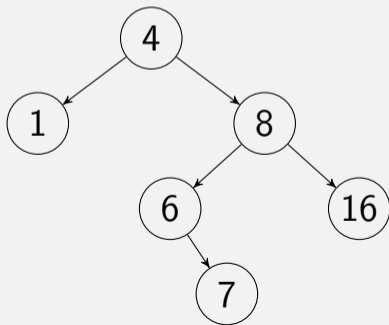
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# Repetition: Binary Trees, Inserting a Key

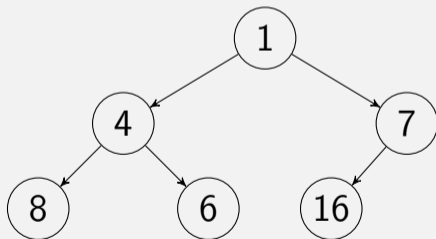
## Binary Search Trees

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# Repetition: Binary Trees, Deleting a Key

## Binary Search Trees

- Replace key  $k$  by symmetric successor  $n$ .
- Careful: What about right child of  $n$ ?

## MinHeap

- Replace key by last element of the array.
- Restore Heap-Condition: `siftDown` or `siftUp`.

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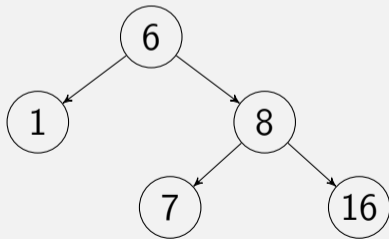
- Replace key by last element of the array.
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**Exercise:** Delete 4 from Example Tree/Heap.

# Repetition: Binary Trees, Deleting a Key

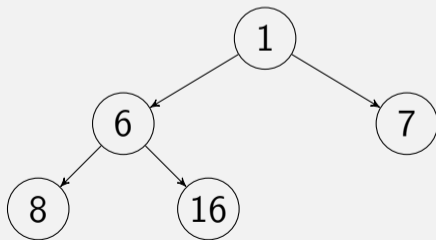
## Binary Search Trees

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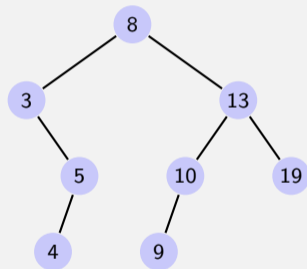
## MinHeap

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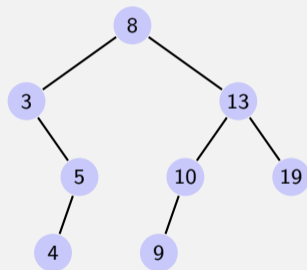
# Traversal possibilities

- preorder:  $v$ , then  $T_{\text{left}}(v)$ , then  $T_{\text{right}}(v)$ .
- postorder:  $T_{\text{left}}(v)$ , then  $T_{\text{right}}(v)$ , then  $v$ .
- inorder:  $T_{\text{left}}(v)$ , then  $v$ , then  $T_{\text{right}}(v)$ .



# Traversal possibilities

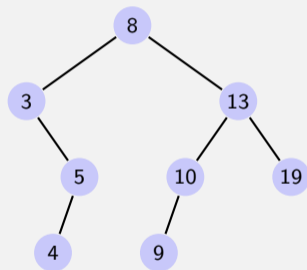
- preorder:  $v$ , then  $T_{\text{left}}(v)$ , then  $T_{\text{right}}(v)$ .  
8, 3, 5, 4, 13, 10, 9, 19
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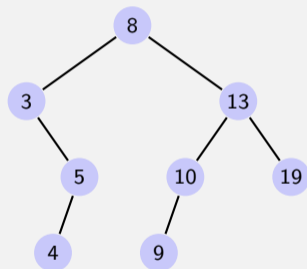
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8, 3, 5, 4, 13, 10, 9, 19
- postorder:  $T_{\text{left}}(v)$ , then  $T_{\text{right}}(v)$ , then  $v$ .  
4, 5, 3, 9, 10, 19, 13, 8
- inorder:  $T_{\text{left}}(v)$ , then  $v$ , then  $T_{\text{right}}(v)$ .



# Traversal possibilities

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8, 3, 5, 4, 13, 10, 9, 19
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4, 5, 3, 9, 10, 19, 13, 8
- inorder:  $T_{\text{left}}(v)$ , then  $v$ , then  $T_{\text{right}}(v)$ .  
3, 4, 5, 8, 9, 10, 13, 19



# Quiz

Draw a binary search tree each that represents the following traversals. Is the tree unique?

|           |                 |
|-----------|-----------------|
| inorder   | 1 2 3 4 5 6 7 8 |
| preorder  | 4 3 1 2 8 6 5 7 |
| postorder | 1 3 2 5 6 8 7 4 |

Provide for each order a sequence of numbers from  $\{1, \dots, 4\}$  such that it cannot result from a valid binary search tree

# Answers

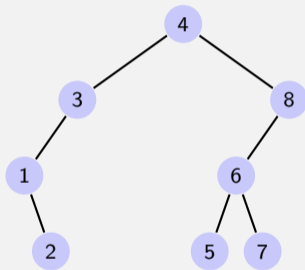
inorder: any binary search tree with numbers  $\{1, \dots, 8\}$  is valid.

The tree is not unique

There is no search tree for any non-sorted sequence. Counterexample  
1 2 4 3

# Answers

preorder 4 3 1 2 8 6 5 7

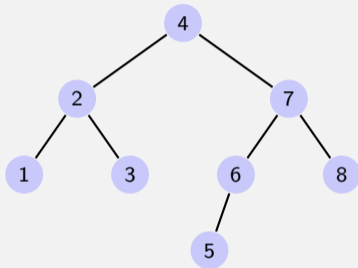


Tree is unique

It must hold recursively that first there is a group of numbers with lower and then with higher number than the first value. Counterexample: 3 1 4 2

# Answers

postorder 1 3 2 5 6 8 7 4

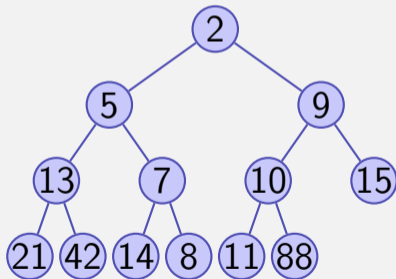


Tree is unique

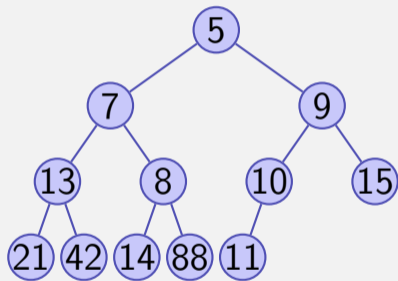
Construction here: <https://www.techiedelight.com/build-binary-search-tree-from-postorder-sequence/>, similar argument as before, but backwards. Counterexample 4 2 1 3

# Heap

On the following Min-Heap, perform an extract-min operation, including re-establishing the heap-condition, as shown in class. What does the heap look like after the operation?



# Solution

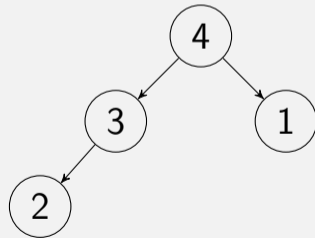
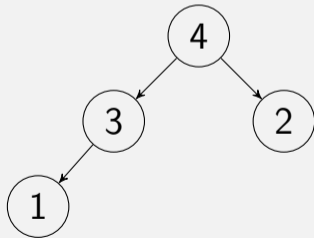
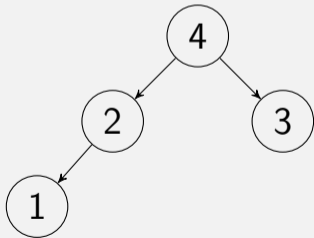




# Quiz: Number of MaxHeaps on $n$ keys

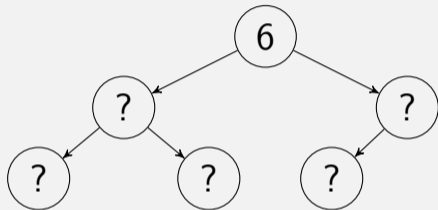
Let  $N(n)$  denote the number of distinct Max-Heaps which can be built from all the keys  $1, 2, \dots, n$ . For example we have  $N(1) = 1$ ,  $N(2) = 1$ ,  $N(3) = 2$ ,  $N(4) = 3$  und  $N(5) = 8$ .

Find the values  $N(6)$  and  $N(7)$ .



# Number of MaxHeaps on $n$ distinct keys

A MaxHeap containing the elements 1, 2, 3, 4, 5, 6 has the structure:



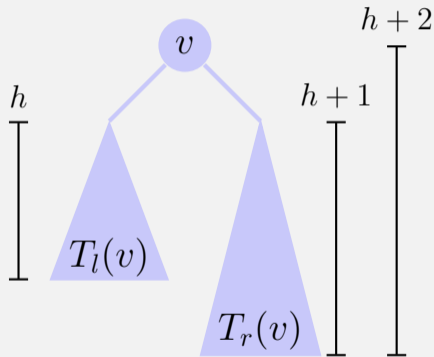
Number of combinations to choose elements for the left subtree:  $\binom{5}{3}$ .

$$\Rightarrow N(6) = \binom{5}{3} \cdot N(3) \cdot N(2) = 10 \cdot 2 \cdot 1 = 20.$$

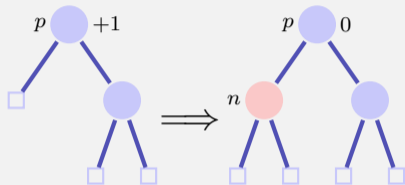
$$\text{and } N(7) = \binom{6}{3} \cdot N(3) \cdot N(3) = 20 \cdot 2 \cdot 2 = 80.$$

# AVL Condition

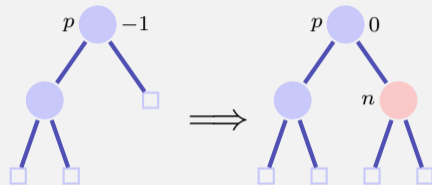
*AVL Condition: for each node  $v$  of a tree  $\text{bal}(v) \in \{-1, 0, 1\}$*



# Balance at Insertion Point



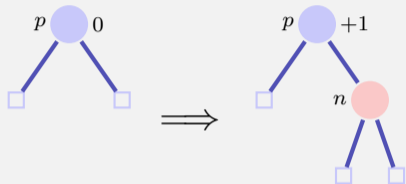
case 1:  $\text{bal}(p) = +1$



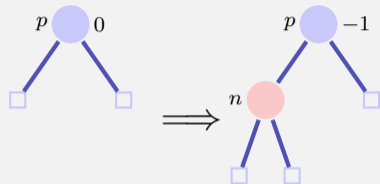
case 2:  $\text{bal}(p) = -1$

Finished in both cases because the subtree height did not change

# Balance at Insertion Point



case 3.1:  $\text{bal}(p) = 0$  right



case 3.2:  $\text{bal}(p) = 0$ , left

Not finished in both case. Call of `upin(p)`

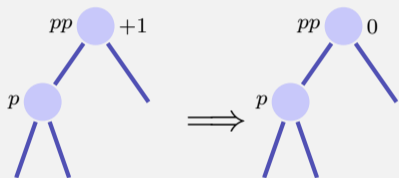
# upin(p) - invariant

When `upin(p)` is called it holds that

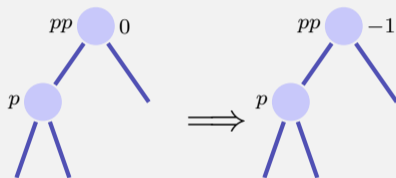
- the subtree from  $p$  is grown and
- $\text{bal}(p) \in \{-1, +1\}$

# upin(p)

Assumption:  $p$  is left son of  $pp^1$



case 1:  $\text{bal}(pp) = +1$ , done.



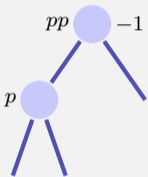
case 2:  $\text{bal}(pp) = 0$ , **upin(pp)**

In both cases the AVL-Condition holds for the subtree from  $pp$

<sup>1</sup>If  $p$  is a right son: symmetric cases with exchange of  $+1$  and  $-1$

# upin(p)

Assumption:  $p$  is left son of  $pp$



case 3:  $\text{bal}(pp) = -1,$

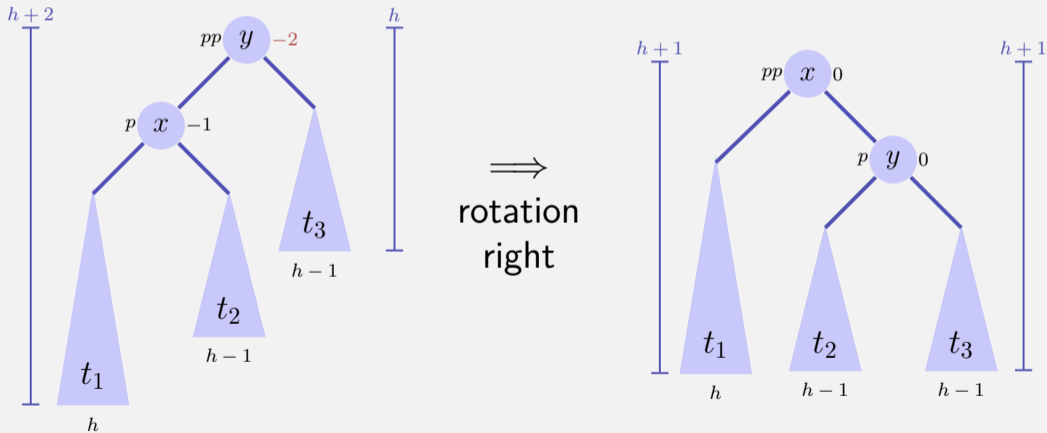
This case is problematic: adding  $n$  to the subtree from  $pp$  has violated the AVL-condition. Re-balance!

Two cases  $\text{bal}(p) = -1, \text{bal}(p) = +1$



# Rotations

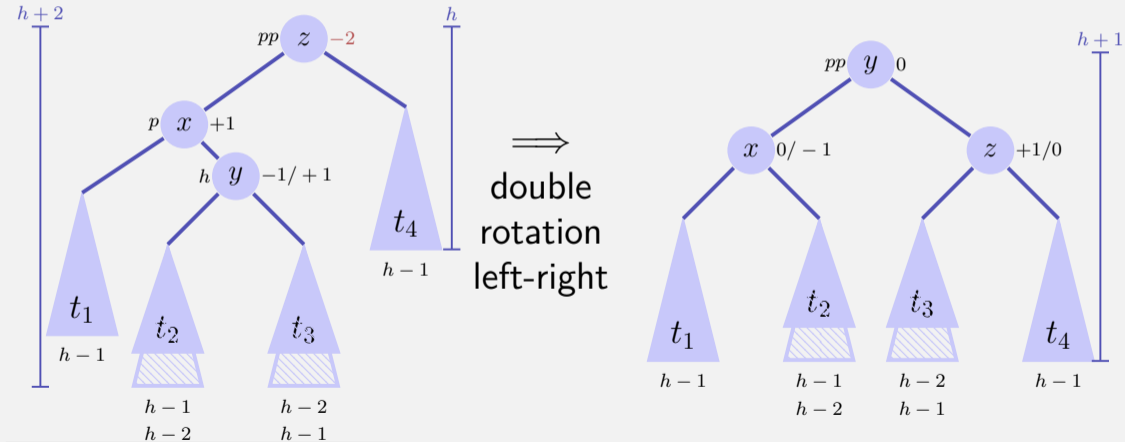
case 1.1  $\text{bal}(p) = -1$ .<sup>2</sup>



<sup>2</sup> $p$  right son:  $\Rightarrow \text{bal}(pp) = \text{bal}(p) = +1$ , left rotation

# Rotations

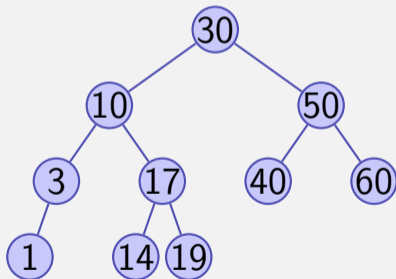
case 1.1  $\text{bal}(p) = -1$ .<sup>3</sup>



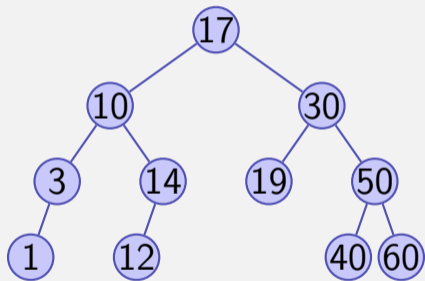
<sup>3</sup> $p$  right son  $\Rightarrow \text{bal}(pp) = +1, \text{bal}(p) = -1$ , double rotation right left

# Quiz

In the following AVL tree, insert key 12 and rebalance (as shown in class). What does the AVL tree look like after the operation that has been shown in class?



# Solution



Questions?