Datenstrukturen und Algorithmen

Exercise 7

FS 2019

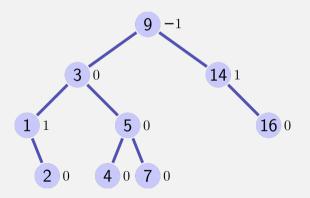
Program of today

1 Feedback of last exercise(s)

2 Repetition theory

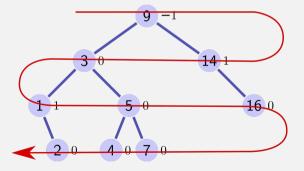
AVL insertion

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- Proof?
- By induction over the height of the tree.
- lacktriangle Hypothesis: Keys at height h and lower are placed in the same place and do not cause insertion.
- lacktriangle Step: Show that the traversal is the same as in the original tree, yields same position. Use AVL property of T to show that there will not be a height difference bigger than 1, and therefore no rotation.

2. Repetition theory

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Longest ascending Sequence in matrix

Given $n \times m$ matrix A:

9	27	42	41	48
35	39	8	3	5
12	49	2	38	4
15	47	29	28	6
19	1	25	33	10

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Wanted longest ascending sequence:

4, 6, 28, 29, 47, 49

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 - In T[x][y] is the length of the longest ascending sequence that ends in A[x][y]
 - In S[x][y] are the coordinates of the predecessor in ascending sequence (if exists)

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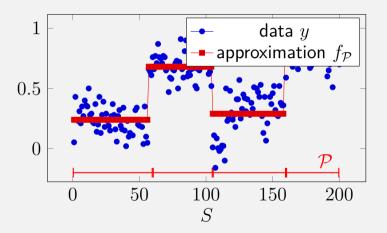
- In which order can entries be computed so that values needed for each entry have been determined in previous steps?
- Start with smallest element in A and so on. (Means that one has to sort A)
- Arbitrary order, if entry is already computed skip it otherwise compute for smaller neighbor recursively.

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■ How can the final solution be extracted once the table has been filled?

Extracting the solution

- How can the final solution be extracted once the table has been filled?
 - Consider all entries to find one with a longest sequence.
 From there, we can reconstruct the solution by following the corresponding predecessors.



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- Goal: find the partition $\hat{\mathcal{P}}$ such that $H_{\gamma,y}(\hat{\mathcal{P}})$ is minimal
- Utilize: efficient computation of the mean using prefix sums (exercise 1): $\mu_I = \frac{1}{|I|} \sum_{i \in I} y_i$

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- **Goal:** find the partition $\hat{\mathcal{P}}$ such that $H_{\gamma,y}(\hat{\mathcal{P}})$ is minimal
- **Dynamic programming**: definition of the table, computation of an entry, calculation order, extracting solution
- Utilize*: $H_{\gamma,y}(\mathcal{P} \cup \{[l,r)\}) = H_{\gamma,y}(\mathcal{P}) + \gamma + e_{[l,r)}$

^{*}Assumption: $\mathcal{P} \cup \{[l,r)\}$ is a partition

Questions?